



Adversarial and Non-Adversarial Risk Analysis for Security over Multiple Sites

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Security Economics: Socio economics meets security

Fare evasion

- Solving only for standard evaders

- Solving only for colluders

- Solving both problems simultaneously

Pickpocketing

Fare evasion and pickpocketing over multiple stations



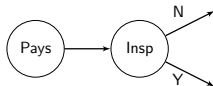
- ▶ One station
- ▶ Two threats.
 - ▶ Fare evasion
 - ▶ Pickpocketing by a team.
 - ▶ Both threats simultaneously.
- ▶ Extension to more than one station.

1. Fare evasion

- ▶ Two types of evaders:
 - ▶ Standard (standard random process).
 - ▶ Colluders (ARA; explicitly modelling intentionality).
- ▶ Five countermeasures.
 - ▶ Inspectors (preventive/recovery), (x_1, c_1) .
 - ▶ Security guards (*bouncers*), usually outsourced (preventive), (x_2, c_2) .
 - ▶ Guards, working solo or in pairs (preventive/recovery), (x_3, c_3) .
 - ▶ Automatic access doors (preventive), $(x_4 \leq n_4, c_4)$.
 - ▶ Metro officers (preventive), $(x_5, \text{negotiation})$.
- ▶ In general, the more resources, the less fare evasion will be. Also, the more inspectors, the more customers checked and, possibly, the more fines collected.

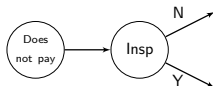
Three types of customers

1. Civic customers.



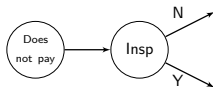
+ Metro	- Metro
Ticket	
Ticket	Annoyed customer

2. Standard fare evaders.



+ Metro	- Metro
	Ticket
Fine	

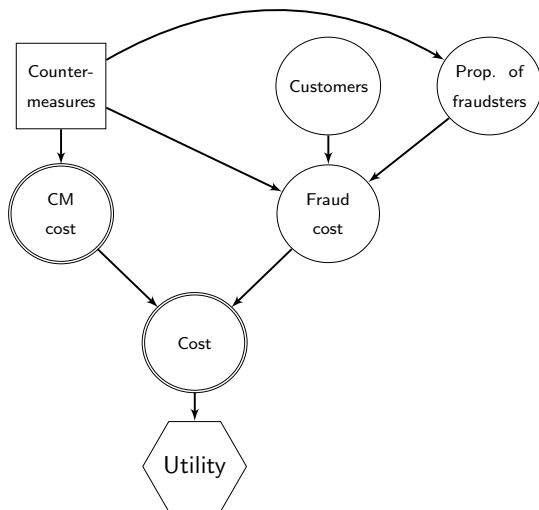
3. Colluding fare evaders.



+ Metro	- Metro	+ Customer	- Customer
	Ticket	Ticket	Cost
Fine			Cost, fine

Solving only for standard evaders

- ▶ This is a 'standard' risk management problem



- ▶ “Club” entailing M operations over incumbent planning period.
 - ▶ They see security investments x (Sequential Defend-Attack).
 - ▶ They decide proportion r of fare evasion
 - ▶ Actual proportion r' depends also on $(x_1, x_2, x_3, x_4, x_5)$.
 - ▶ Operational costs, including preparation costs c_e

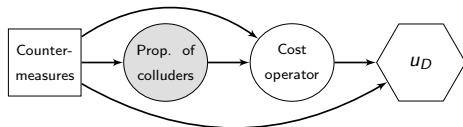
$$c = t(M_2 - M_1) - fM_3 - rc_eM.$$

- ▶ M_1 evaders pay (abortion).
- ▶ M_2 not pay, not caught.
- ▶ M_3 not pay, caught.

$$(M_1, M_2, M_3) \sim \mathcal{M}(M; (1 - r'), r'(1 - q_A(x_1)), r'q_A(x_1)).$$

- ▶ Their utility (risk prone)

$$u_C(c - rc_eM).$$



- ▶ Relevant revenues and costs

Concept	Revenue
Security costs	$-(c_1x_1 + c_2x_2 + c_3x_3 + c_4x_4 + x_5)$
Tickets lost	$-tM_2$
Fines won	fM_3

- ▶ Total increase in outcome is

$$c_D = fM_3 - tM_2 - (c_1x_1 + c_2x_2 + c_3x_3 + c_4x_4 + x_5).$$

- ▶ $u_D \rightarrow$ utility, $h(r|x)$ models their beliefs over r given x , then

$$\psi(x) = \int \left[\sum_{M_1, M_2, M_3} p_{M_1 M_2 M_3 x} u_D \left(fM_3 - tM_2 - \sum_{i=1}^4 c_i x_i - x_5 \right) \right] \times h(r|x) dr,$$

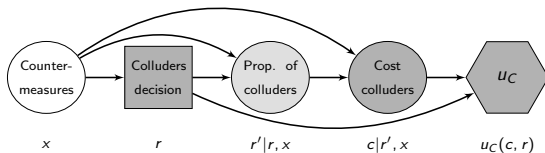
$$p_{M_1 M_2 M_3 x} = \Pr(M_i \text{ type } i \text{ colluders} | x \text{ invested}), \sum_{i=1}^3 M_i = M.$$

- ▶ They must solve

$$\max_{x \in \mathcal{B}_1} \psi(x_1, x_2, x_3, x_4, x_5),$$

but $h(r|x)$ not directly available (\rightarrow Attacker's problem).

Attacker's problem



- ▶ x , security investment by Operator. We may consider $p_A(x)$ (Attacker's beliefs over x), although not necessary (is seen).
- ▶ r , decision made by Attacker.
- ▶ r' , effective fare evasion proportion.
 - ▶ One possible model $r' = r(1 - s(x_1))$, $s(x_1)$ proportion of evasion abortions.
 - ▶ $p_A(s(x_1)) = p_A(s|x_1)$ induces $p_A(r'|r, x)$.
- ▶ c , global costs of evasion operation.
- ▶ u_C , utility over consequences $u_C(c - rc_e M)$.

Solving Attacker's problem



1. Integrate out uncertainty over c , getting expected utility

$$\psi_A(r', r, x) = \int \left[\sum_{M_1, M_2, M_3} p_{M_1 M_2 M_3 x} \times u_C(t(M_2 - M_1) - fM_3 - rc_e M) \right] \times g_A(q_A | x_1) dq_A.$$

$g_A(q_A | x_1)$ density over $q_A | x_1$, inducing $p_A(c | r', x)$.

2. Integrate out uncertainty over r' , obtaining expected utility

$$\psi_A(r, x) = \int \psi_A(r', r, x) p_A(s | x_1) ds.$$

3. Find Attacker's optimal strategy

$$r(x) = \arg \max_r \psi_A(r, x).$$

Simulation scheme for estimating $h(r|x)$



- Uncertainty about $u_C(\cdot)$, $g_A(q_A|\cdot)$, $p_A(s|\cdot)$, modelled through $U_C(\cdot)$, $G_A(q_A|\cdot)$, $P_A(s|\cdot)$, has to be propagated.

For each x

For $i = 1$ to K

Sample $U_C^i, G_A^i(q_A|\cdot), P_A^i(s|\cdot)$. Compute

$$\psi_A^i(r', r, x) =$$

$$\int \left[\sum_{M_1, M_2, M_3} p_{M_1 M_2 M_3 x} U_C^i(t(M_2 - M_1) - fM_3 - rc_e M) \times G_A^i(q_A|x_1) \right] dq_A.$$

Compute

$$\psi_A^i(r, x) = \int \psi_A^i(r', r, x) P_A^i(s|x_1) ds.$$

Compute random optimal alternative

$$R^i = \operatorname{argmax}_r \psi_A^i(r, x).$$

Approximate $p_A(R(x) \leq r) \approx \#\{1 \leq i \leq K : R^i \leq r\}/K$.

Typical assumptions

- ▶ Colluders risk prone in benefits $\rightarrow u_C$ strategically equivalent

$$u_C(c) = \exp(k_C c), k_C > 0.$$

- ▶ A random utility model could be

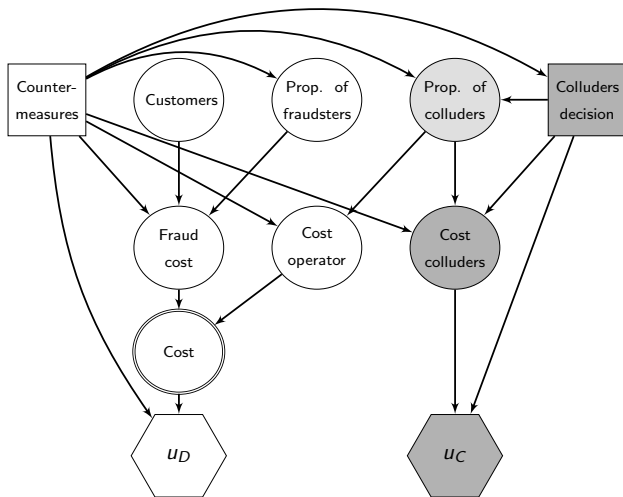
$$U_C(c) = \exp(k_C c), k_C \sim \mathcal{U}(0, K_C).$$

- ▶ Evaders proportion $s \sim \mathcal{B}e(\alpha, \beta)$. Dirichlet process with base $\mathcal{B}e(\alpha, \beta)$ for P_A

$$P_A \sim \mathcal{D}\mathcal{P}(\mathcal{B}e(\alpha, \beta), \delta_1).$$

- ▶ If we consider $r' > r$, we could use an error model $r' = r + s$, s described by $p_A(s)$ and $P_A \sim \mathcal{D}\mathcal{P}(p_A, \delta_2)$.
- ▶ Proportion of inspections $q_A(x_1) \sim \mathcal{B}e(\alpha, \beta)$ with $\frac{\alpha}{\alpha + \beta} = \delta_{x_1}$ and small variance.
 - ▶ Then, $G_A \sim \mathcal{D}\mathcal{P}(\mathcal{B}e(\alpha, \beta), \delta_3)$.

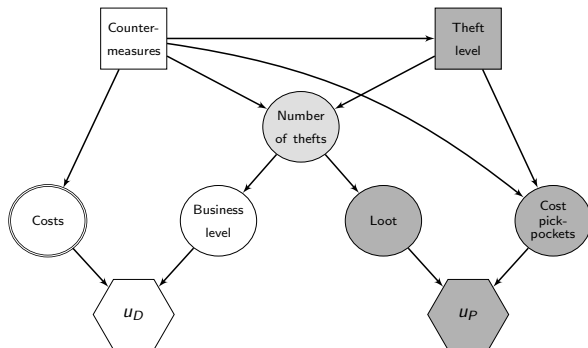
Solving the problem when both evaders are present



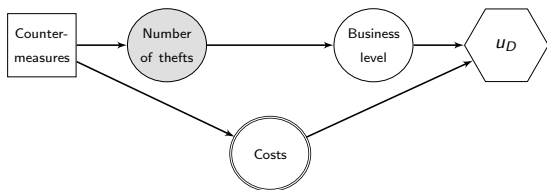
2. Pickpocketing

Description of problem

- ▶ Four countermeasures.
 - ▶ Patrols (preventive/recovery), (y_1, d_1) .
 - ▶ Cameras (preventive), (y_2, d_2) .
 - ▶ Guards (preventive/recovery), (x_3, c_3) .
 - ▶ Public awareness plans (preventive), (y_3) .



Defender's problem



- ▶ Operator invests y_1, y_2, x_3 (units) and y_3 (in the plan).
- ▶ Faces a delinquency level.
- ▶ Sees a decrease in business.
- ▶ Gets her utility (depends on business level and operator costs).

Defender's problem (cont.)



- ▶ Security costs

$$c'_D(y_1, y_2, x_3, y_3) = d_1 y_1 + d_2 y_2 + c_3 x_3 + y_3.$$

- ▶ b , business level, T theft level, $u_D(c'_D, b)$ Defender's utility

$$\max_{y \in \mathcal{B}_2} \iint u_D(c'_D, b) p(b|T) p(T|y) dT dy.$$

- ▶ $u_D(c'_D, b)$ includes costs and reduction in business level

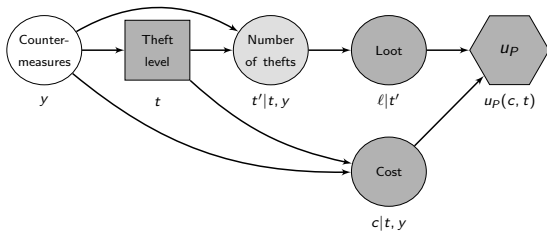
$$c'_D + (b_0 - b).$$

- ▶ Operator typically risk averse with respect to costs

$$u_D(c'_D, b) = -\exp(k'_D \cdot [c'_D + (b_0 - b)]), \quad k'_D > 0.$$

- ▶ To assess $p(T|y)$ → Attacker's problem.

Attacker's problem



- ▶ See Defender's investment (y_1, y_2, x_3, y_3) .
- ▶ Decide on target theft level T .
- ▶ Implement actual number of theft operations, $T' = \tau T$.
- ▶ Costs (of implementing) their actions is $c_p T$.

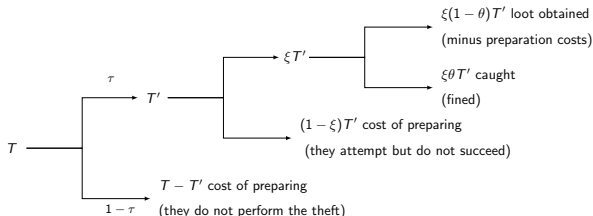
Attacker's problem (cont.)



- ▶ Face operational costs.
 - ▶ With prob. $(1 - \xi)$, unsuccessful attempt. No consequences.
 - ▶ With prob. $\xi\theta$, succeed but caught. Fine g .
 - ▶ With prob. $\xi(1 - \theta)$, succeed and not caught. Loot L .
- ▶ Total cost/benefit balance

$$c = [-c_p \times T_1] - [(g + c_p) \times T_2] + [(L - c_p) \times T_3] = -c_p T - gT_2 + LT_3.$$

- ▶ Get utility $u_P(-c_p T - gT_2 + LT_3)$.



For each y

For $i = 1$ to K

Sample $U_P^i, P_A^i(\tau|\cdot), P_A^i(\xi|\cdot), P_A^i(\theta|\cdot)$.

Compute

$$\psi_P^i(t', t, y) = \iint \left[\sum_{T_1, T_2, T_3} P_{T_1 T_2 T_3 y} \int U_P^i(-c_P T - g T_2 + L T_3) dU_P \right] \\ \times P_A^i(\xi|y_1, x_3, y_3) P_A^i(\theta|y_1, x_3) d\xi d\theta.$$

Compute

$$\psi_P^i = \int \psi_P^i(t', t, y) P_A^i(\tau|y_1, x_3) d\tau.$$

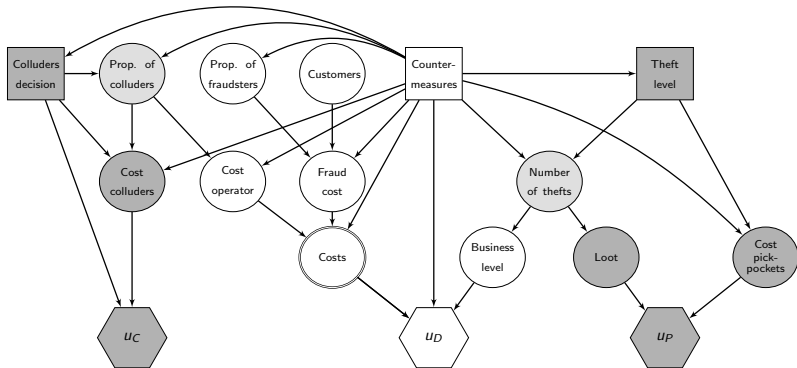
Compute (and register) the optimal alternative

$$T^i = \operatorname{argmax}_t \psi_P^i(t, y).$$

Approximate $p_A(T(y) \leq t) \approx \#\{1 \leq i \leq K : T^i \leq t\}/K$.

3. Fare evasion and pickpocketing over multiple stations

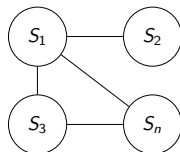
Joint influence diagram



Multiple sites problem



- ▶ Colluders and pickpockets do not make common cause.
- ▶ We can solve their problems separately.
- ▶ A network of n interconnected stations



- ▶ For each station a model like above is applicable, with mobile resources subject to global and specific budget and mobility constraints.



- ▶ Fare evaders
 - ▶ If too many security measures in entering station i , move to adjacent station i' .
 - ▶ If inspectors in intermediate station k , an alternative route.
- ▶ Pickpockets
 - ▶ If too many security measures in station i , move to adjacent station i' .
- ▶ Some personnel (inspectors, patrols, guards) are mobile.

- ▶ A DSS is being currently devised to help decision makers.
- ▶ Upon perceived low-level threats, authorities tend to underestimate risk.
 - ▶ Attackers see a breach in security (more attackers).
 - ▶ Great impact can be caused even with low-profile attacks.
 - ▶ Low-cost preventive measures and well-trained personnel could deter attackers or minimize their number.
- ▶ Under scenario of high probability of attack.
 - ▶ Authorities tend to invest on expensive (sometimes sensationalist and ineffective) measures.
 - ▶ Set up security and mobility protocols for personnel increase their efficiency.

- ▶ For fare evasion, a mixed non adversarial - adversarial problem has been tackled.
- ▶ For pickpockets, not only direct economic impact considered (also image costs).
- ▶ General model over multiple sites has been devised.
 - ▶ Resources are constrained by budget and mobility.
 - ▶ Some countermeasures have to be shared for both threats.